

## Concurrency Theory

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### Sheet 6

**Due: Monday, 2025-12-08**

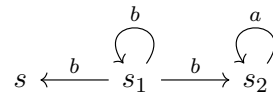
#### Exercise 6.1

Use two separate formulas to describe the properties in Hennessey-Milner Logic with recursion.

- a) There is an infinite  $a$ -labelled computation path in which all states have an outgoing  $b$ -labelled transition.
- b) Every  $b$ -labelled computation path leads to a state from which an  $a$ -labelled transition is not possible.

#### Exercise 6.2

Consider the LTS



Compute all fixpoints of the functions

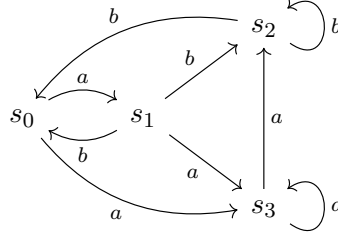
- a)  $\llbracket \langle a \rangle tt \vee [b]X \rrbracket$
- b)  $\llbracket \langle a \rangle tt \vee ([b]X \wedge \langle b \rangle tt) \rrbracket$

#### Exercise 6.3

Given a LTS  $L = (S, Act, \rightarrow)$ , show that  $\llbracket F \rrbracket : 2^S \rightarrow 2^S$  is a monotonic function over the complete lattice  $(2^S, \subseteq)$ , for all formulas  $F$ , expressible in HML with recursion.

**Exercise 6.4**

Consider the LTS



- Compute  $\llbracket \langle b \rangle [a] tt \wedge \langle b \rangle [b] X \rrbracket (\{s_0, s_2\})$
- Compute the set of processes satisfying the property

$$X \stackrel{\min}{=} \langle b \rangle \langle a \rangle tt \vee \langle b \rangle [b] X$$

- Compute the sets of processes satisfying the mutual recursive equational system

$$A \stackrel{\max}{=} [a] B$$

$$B \stackrel{\max}{=} \langle a \rangle C \wedge [b] B$$

$$C \stackrel{\max}{=} [b] B$$