Concurrency Theory

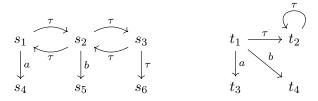
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Sheet 3

Due: Monday, 2025-11-17

Exercise 3.1

Consider the two LTS below:



Show that $s_1 \approx t_1$ by constructing a weak bisimulation.

Exercise 3.2

Prove or disprove the statements below:

 $a.\tau.nil \sim \tau.a.nil$

 $a.\tau.nil \approx \tau.a.nil$

Exercise 3.3

Prove that $P + \tau . P \approx \tau . P$ for all CCS processes P.

Exercise 3.4

In the weak bisimulation game the attacker is only allowed to use $\stackrel{a}{\rightarrow}$ moves, but the defender is allowed to use $\stackrel{a}{\Rightarrow}$ moves. Give an argument why letting the attacker use $\stackrel{a}{\Rightarrow}$ moves would not make the attacker more powerful.

Exercise 3.5

Define two CCS processes A and B such that

- LTS(A) has an infinite number of states,
- LTS(B) has a finite number of states
- $A \sim B$ is false (no proof needed)
- $A \approx B$ is true (no proof needed)

If you have questions, please post a message in the dedicated chat.

Bonus: prove that $A \approx B$ is true, but $A \sim B$ is false.

Exercise 3.6 (Bonus)

We want to model a simple communication protocol between a sender and a reciever. The sender accepts data and the reciever delivers the data. We want to implement this specification via a medium that transmits the data between sender and reciever.

Consider the following CCS definition for the specification and implementation of this communication protocol:

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\label{eq:protocolSpec} \begin{split} &\operatorname{ProtocolSpec} \\ &\operatorname{Protocol} \doteq (\operatorname{Send} \parallel \operatorname{Med} \parallel \operatorname{Rec}) \setminus \{\operatorname{send}, \operatorname{err}, \operatorname{trans}, \operatorname{ack}\} \\ &\operatorname{Send} \doteq \operatorname{acc.Sending} \\ &\operatorname{Med} \doteq \operatorname{send.Trans} \\ &\operatorname{Rec} \doteq \operatorname{trans.Del} \end{split}
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Send = acc.Sending Med = send.Frans Rec = trans.Del.

Sending $\stackrel{.}{=}$ $\overline{\text{send}}$.Wait $\overline{\text{Trans}} \stackrel{.}{=} \tau$.Err + $\overline{\text{trans}}$.Med Del $\stackrel{.}{=}$ $\overline{\text{del}}$.Ack

Wait $\stackrel{.}{=}$ ack.Send + err.Sending $\overline{\text{Err}} \stackrel{.}{=} \overline{\text{err}}$.Med Ack $\stackrel{.}{=}$ $\overline{\text{ack}}$.Rec

Show that Protocol is an implementation of ProtocolSpec by showing that Protocol \approx ProtocolSpec is true.